Designing Features for Mature Systems: Lessons Learned from Manta

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Legacy?

distributed legacy cloud opensource

Manta 101!



Manta 101: Features

- Highly scalable, distributed object store
- HTTP-based
- Compute as a first class citizen!
 - Manta "jobs"
- Interface feels like a Unix filesystem

```
/jhendricks/stor/myDir/myObj.txt
```



Manta 101: API operations

- mput: upload an object
- mget: download an object
- mrm: remove an object
- mmkdir: create a directory
- mrmdir: remove a directory
- mls: list a directory
- mln: make a 'snaplink' to an object
- mlogin: login to a container with your object

Manta 101: System Primitives

- Objects
 - Stored as flat files
 - Immutable
- Directories
 - Are listable!



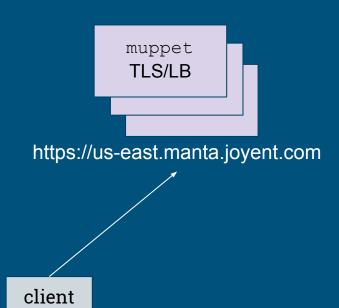
Manta 101: Design Constraints

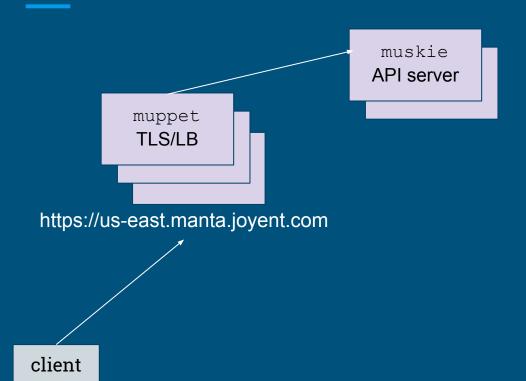
- CP system
 - When faced with a network partition, Manta chooses consistency over availability.
- Horizontally scalable where possible

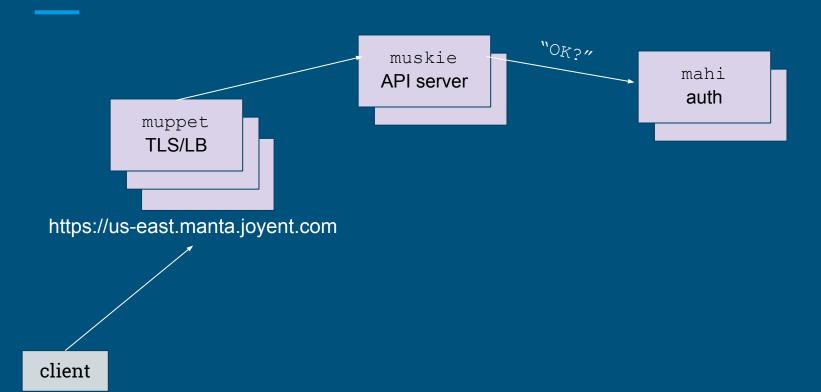


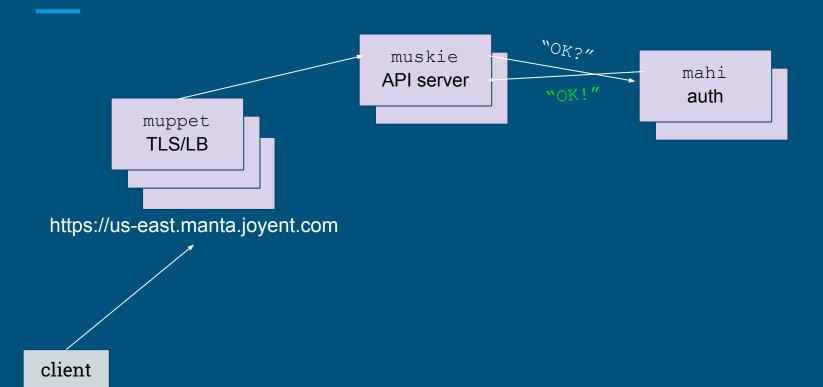
https://us-east.manta.joyent.com





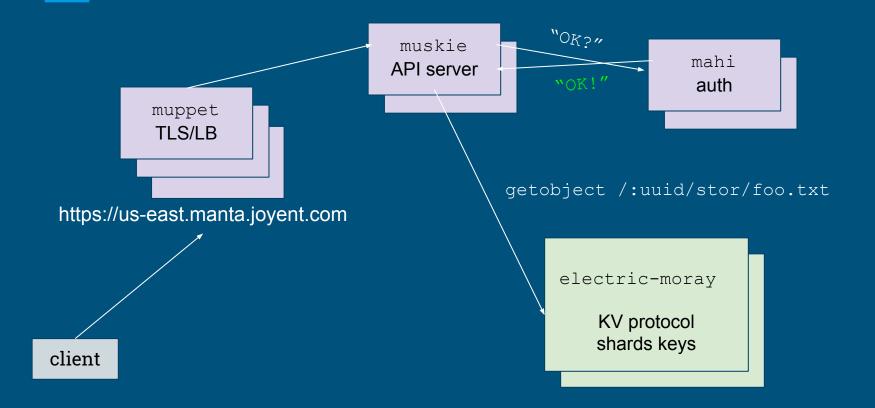






Manta Design: Data Path

- Architectural separation of metadata and storage
- Metadata tier responsible for information about the objects: its name, size, content MD5, who owns it, permissions, where it is stored, etc.
- Storage tier responsible for storing the object
 - Compute jobs also run directly on data on storage
 CNs



electric-moray

KV protocol shards keys

2.moray

KV store

3.moray

KV store

4.moray

KV store

electric-moray

KV protocol
shards keys

hash(dirname("/:uuid/stor/foo.txt"))
=> 3.moray

2.moray

KV store

3.moray

KV store

4.moray

KV store

2.moray electric-moray KV store getobject /:uuid/stor/foo.txt KV protocol shards keys 3.moray KV store => 3.moray 4.moray KV store

electric-moray

KV protocol
shards keys

getobject /:uuid/stor/foo.txt

KV store

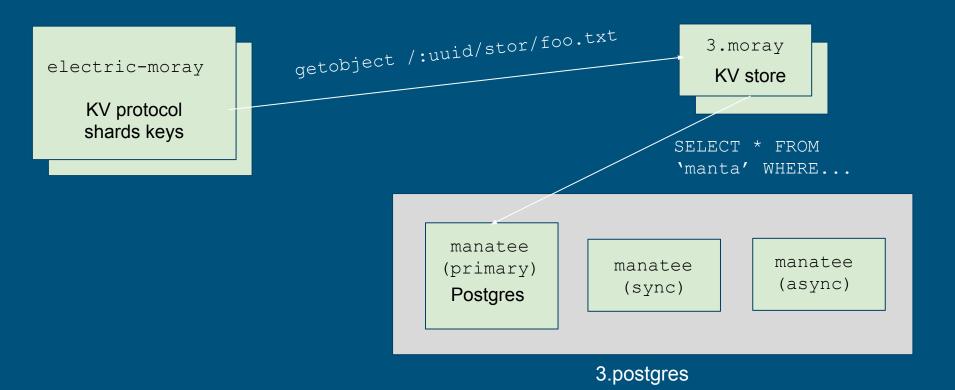
KV store

manatee
(primary)
Postgres

manatee
(sync)

manatee
(async)

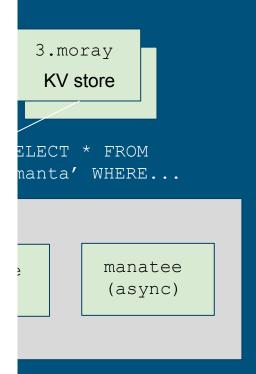
3.postgres



electric-moray

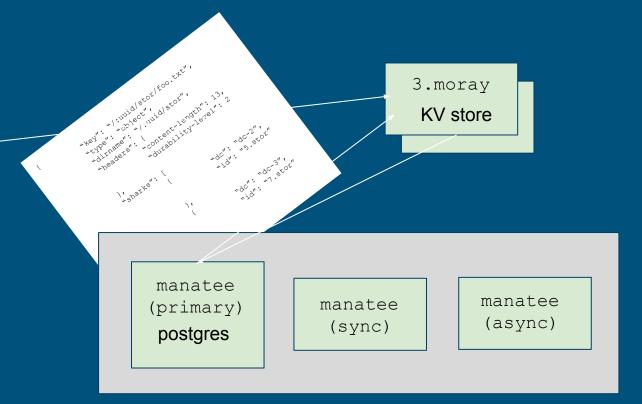
KV protocol shards keys

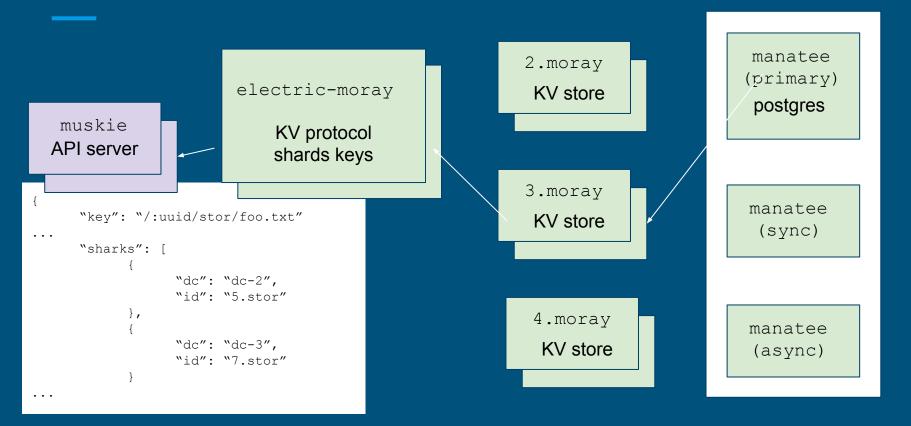
```
"key": "/:uuid/stor/foo.txt",
"type": "object",
"dirname": "/:uuid/stor",
"headers": {
     "content-length": 13,
     "durability-level": 2,
     "content-type": "application/text"
"sharks": [
          "dc": "dc-2",
          "id": "5.stor"
     },
          "dc": "dc-3",
          "id": "7.stor"
```



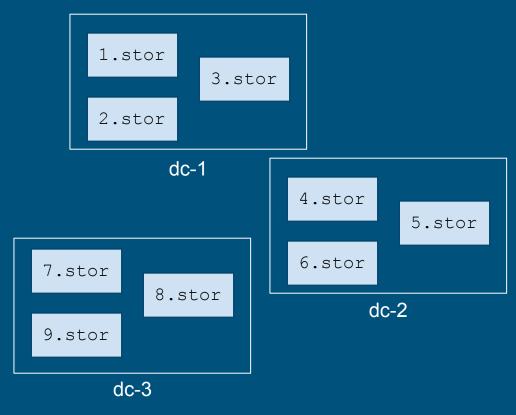
electric-moray

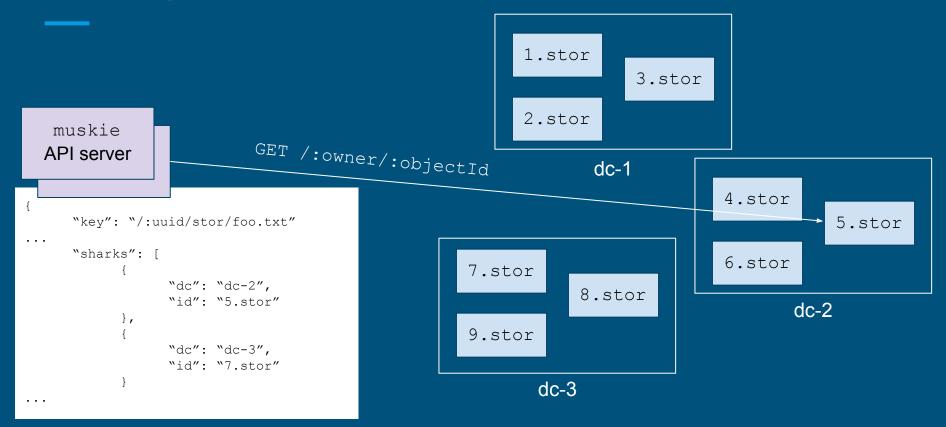
KV protocol shards keys

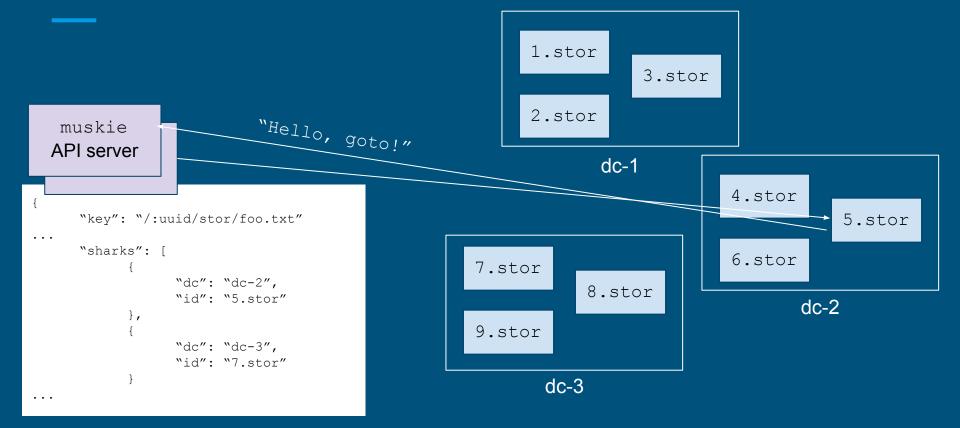




```
muskie
API server
   "key": "/:uuid/stor/foo.txt"
   "sharks": [
               "dc": "dc-2",
               "id": "5.stor"
         } ,
               "dc": "dc-3",
               "id": "7.stor"
```

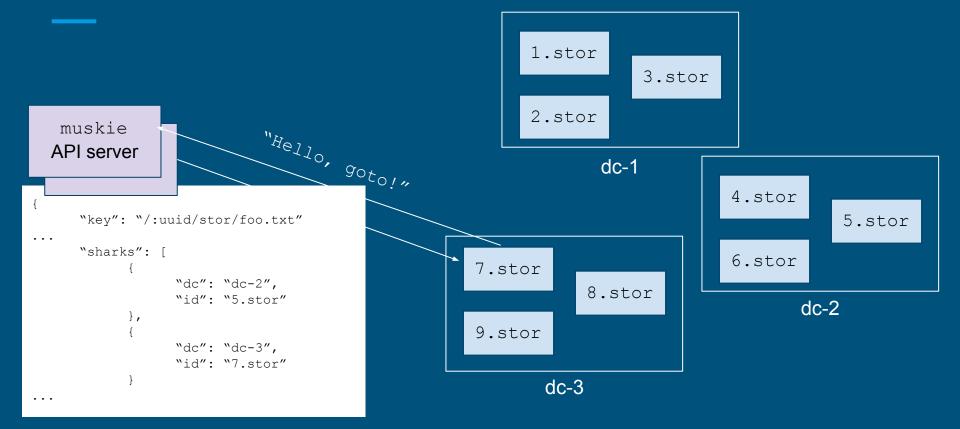






```
1.stor
                                                                   3.stor
                                                      2.stor
 muskie
API server
                                                              dc-1
                                                                             4.stor
   "key": "/:uuid/stor/foo.txt"
                                                                                          5.stor
   "sharks": [
                                                                             6.stor
                                                7.stor
              "dc": "dc-2",
                                                             8.stor
              "id": "5.stor"
                                                                                     dc-2
        },
                                                9.stor
              "dc": "dc-3",
              "id": "7.stor"
                                                        dc-3
```

```
1.stor
                                                                 3.stor
                                                     2.stor
muskie
                 GET /:owner/:objectId
API server
                                                             dc-1
                                                                           4.stor
   "key": "/:uuid/stor/foo.txt"
                                                                                       5.stor
   "sharks": [
                                                                           6.stor
                                               7.stor
             "dc": "dc-2",
                                                           8.stor
             "id": "5.stor"
                                                                                   dc-2
        },
                                               9.stor
             "dc": "dc-3",
             "id": "7.stor"
                                                      dc-3
```



PUT /jhendricks/stor/newObj.txt

- Authorization (mahi)
- Muskie picks sharks to store the object on (spread across at least 2 DCs)
- Muskie streams the data to the sharks
- Muskie updates electric-moray with the new metadata record

Manta Design: Compute

- Objects can be large
- Want to avoid copying data between servers
- To make objects possible to compute on with existing software, Manta's abstraction for an object is a flat file on its storage node.
 - Can run programs directly on these files
- Instead of copying data to run jobs on, move jobs to where the data lives!

\$ wc /jhendricks/stor/foo.txt

muskie API server client

\$ cat foo.txt Hello, goto!

5.stor

\$ wc /jhendricks/stor/foo.txt



\$ wc /jhendricks/stor/foo.txt



Multipart Uploads!



Multipart Uploads: First Principles

- Upload an object in "parts"
- "Commit" the object when all parts are uploaded
 - Creates a new object, indistinguishable from normal objects created via PUT path
- Use cases?
 - Large files
 - Intermittent network connections
 - Streaming data from a source

Multipart Uploads: Design Questions

- How to identify multipart uploads?
- Where to store parts?
- What does a "commit" of an MPU look like?

Multipart Uploads: Design Considerations

- How can we ensure parts are listed easily?
- How can we list all multipart uploads?
- What happens if multiple clients operate on the same MPU?
- What happens during failures of Manta components?
- How will parts be cleaned up after commit?
- ...Can I cancel MPUs, too?

Multipart Upload Design

- How to identify MPUs?
 - UUID per MPU
- How to store parts?
 - Obvious answer: Use Manta objects!
 - Store parts as objects in a Manta directory
 - One directory per MPU: max of 10,000 parts << number allowed dirents (1 million)
 - Allows clients to list parts easily

Multipart Upload Design

- How to store part directories?
 - o Under a new top-level directory, /:account/uploads (analogous
 to /:account/jobs)
 - But: don't want to limit the number of ongoing MPUs to the number of allowed dirents
 - Solution: one-level nested "prefix" directories, in which all uploads starting with the same characters have the same parent
 - Allows all MPUs to be listed in as many requests as it takes to list all prefix directories

Example MPU Structure

- Upload ID: eaff0760-9b17-4fb7-b7c4-f2de818681f3
- Parts directory

```
/jordan/uploads/eaf/eaff0760-9b17-4fb7-b7c4-f2de818681f3
```

Example parts

```
/jordan/uploads/eaf/eaff0760-9b17-4fb7-b7c4-f2de818681f3/0
/jordan/uploads/eaf/eaff0760-9b17-4fb7-b7c4-f2de818681f3/1
/jordan/uploads/eaf/eaff0760-9b17-4fb7-b7c4-f2de818681f3/2
```

Multipart Upload Design: Commits

- Design constraints for mpu-commit operation?
 - Must be idempotent
 - Must be atomic
- What steps need to happen in Manta architecture?
 - Metadata layer: Insert an object record for the target object
 - Storage layer: Create the object on disk from its parts, on the appropriate storage nodes

Commit: Storage Layer Implementation

Constraints

- Cannot append or mutate parts on an existing object
- Would like to avoid copying data over network.

Design

- Co-location of parts on the storage nodes the final target object will live on
- Create object on disk from parts locally
- New operation on mako (storage node service):
 mako-finalize: requires array of part etags

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
```

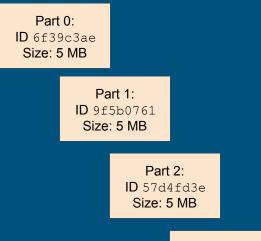
Part 0: ID 6f39c3ae Size: 5 MB

Part 1: ID 9f5b0761 Size: 5 MB

> Part 2: ID 57d4fd3e Size: 5 MB

> > Part 3: ID 5002e70d Size: 1 MB

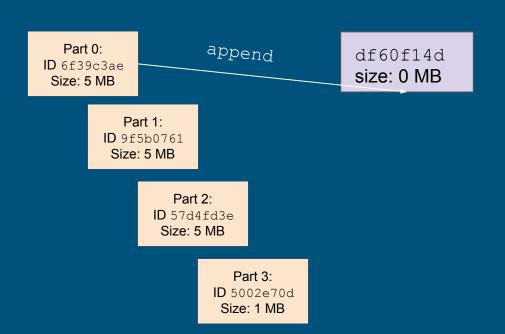
```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
```



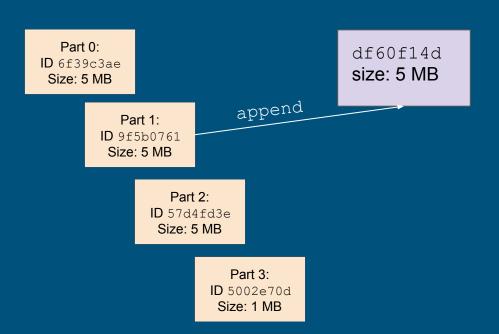
df60f14d

Part 3: ID 5002e70d Size: 1 MB

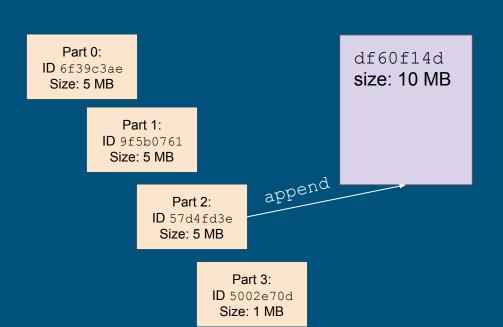
```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
```



```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
```

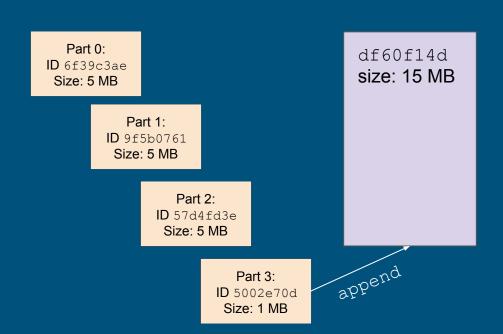


```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
```



3.stor.us-east.joyent.us

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
```



```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
```

Part 0: ID 6f39c3ae Size: 5 MB Part 1: ID 9f5b0761 Size: 5 MB Part 2: **ID** 57d4fd3e Size: 5 MB Part 3: ID 5002e70d Size: 1 MB

df60f14d size: 16 MB

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
```

Part 1: ID 9f5b0761 Size: 5 MB

> Part 2: ID 57d4fd3e Size: 5 MB

> > Part 3: ID 5002e70d Size: 1 MB

df60f14d size: 16 MB

3.stor.us-east.joyent.us

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: 67115618
```

df60f14d size: 16 MB

Part 2: ID 57d4fd3e Size: 5 MB

> Part 3: ID 5002e70d Size: 1 MB

3.stor.us-east.joyent.us

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
```

df60f14d size: 16 MB

Part 3: ID 5002e70d Size: 1 MB

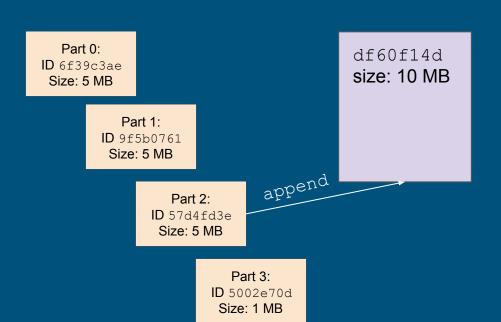
```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
```

df60f14d size: 16 MB

mako-finalize: atomicity & idempotency

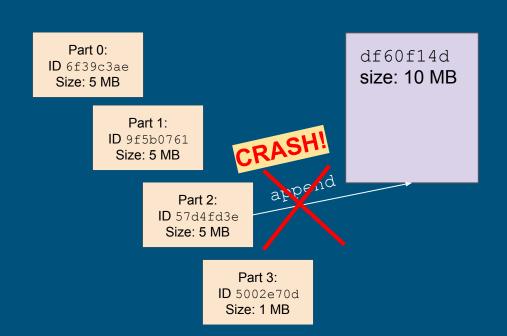
- What happens if make crashes...
 - While writing to the target object file?

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
outputFile: df60f14d
```



3.stor.us-east.joyent.us

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
outputFile: df60f14d
```



mako-finalize: atomicity & idempotency

- What happens if make crashes...
 - While writing the target object file?
 - mako-finalize can safely retry later, because it hasn't removed the parts yet
 - While removing parts?

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
```

Part 0: ID 6f39c3ae Size: 5 MB Part 1: ID 9f5b0761 Size: 5 MB Part 2: **ID** 57d4fd3e Size: 5 MB Part 3: ID 5002e70d Size: 1 MB

df60f14d size: 16 MB

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
```

Part 1: ID 9f5b0761 Size: 5 MB

> Part 2: ID 57d4fd3e Size: 5 MB

> > Part 3: ID 5002e70d Size: 1 MB

df60f14d size: 16 MB

3.stor.us-east.joyent.us

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
```

df60f14d size: 16 MB

Part 2: ID 57d4fd3e Size: 5 MB

> Part 3: ID 5002e70d Size: 1 MB

```
parts: [
   6f39c3ae,
   9f5b0761,
   57d4fd3e,
   5002e70d
targetFile: df60f14d
```

df60f14d size: 16 MB



Part 2: **ID** 57d4fd3e Size: 5 MB

> Part 3: ID 5002e70d Size: 1 MB

mako-finalize: atomicity & idempotency

- What happens if make crashes...
 - While writing the target object file?
 - Can safely retry later, because it hasn't removed the parts yet
 - While removing parts?
 - Can still retry later...
 - Need a way to check that the target object file exists and is correct

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
nbytes:16 MB
```

df60f14d size: 16 MB

Part 2: ID 57d4fd3e Size: 5 MB

> Part 3: ID 5002e70d Size: 1 MB

3.stor.us-east.joyent.us

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
nbytes:16 MB
```

df60f14d size: 16 MB

Part 3: ID 5002e70d Size: 1 MB

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
nbytes:16 MB
```

df60f14d size: 16 MB

3.stor.us-east.joyent.us

mako-finalize ----- Success!

```
parts: [
    6f39c3ae,
    9f5b0761,
    57d4fd3e,
    5002e70d
]
targetFile: df60f14d
nbytes:16 MB
```

df60f14d size: 16 MB

mako-finalize: atomicity & idempotency

- Protections presented thus far prevent problems from the same MPU (exact same set of parts)
- Additional constraints needed at the muskie (REST API) layer
 - Don't want to allow two clients to commit a different set of parts for the same
 MPU
 - Check for conflicts before invoking mako-finalize
 - Other conflicts:
 - Aborts and commits conflict with each other
 - If one client tries to commit an MPU and another tries to abort it, we only want one to win (atomicity)
 - Need to store additional state in the metadata tier

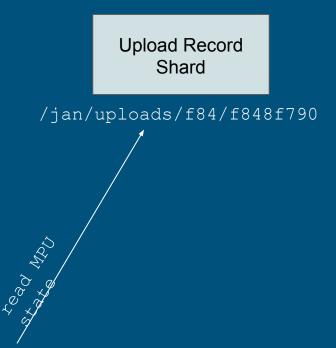
Commits: atomicity & idempotency

- Where do we store the multipart upload state?
- Considerations:
 - Need the state change and target object's visibility in Manta to be an atomic operation
- Suppose state about an MPU was stored in only the metadata record of the parts directory...
 - o Is this atomic?

Upload Record Shard

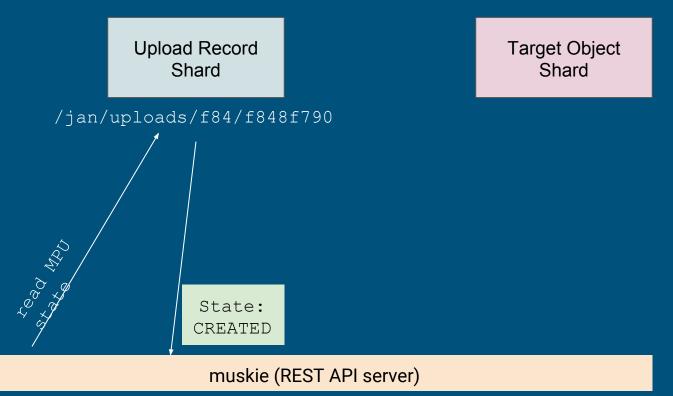
Target Object Shard

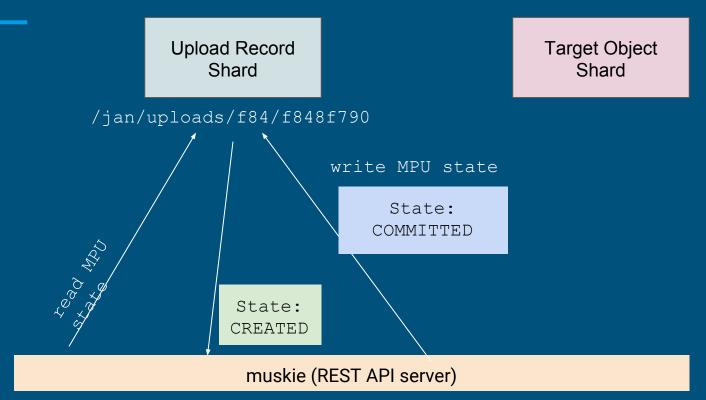
/jan/uploads/f84/f848f790

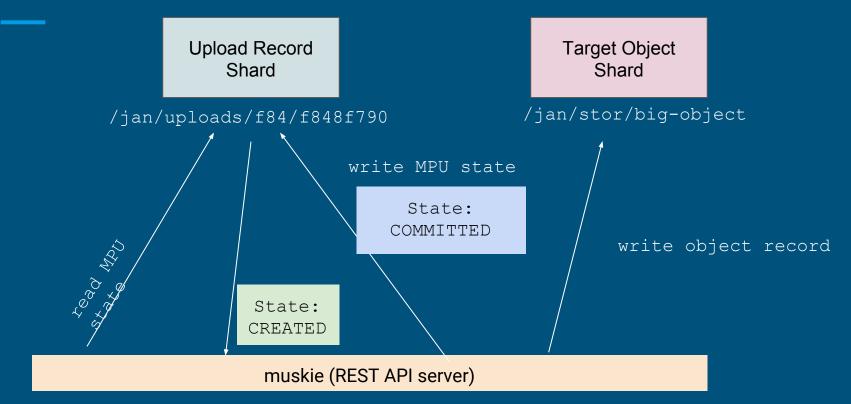


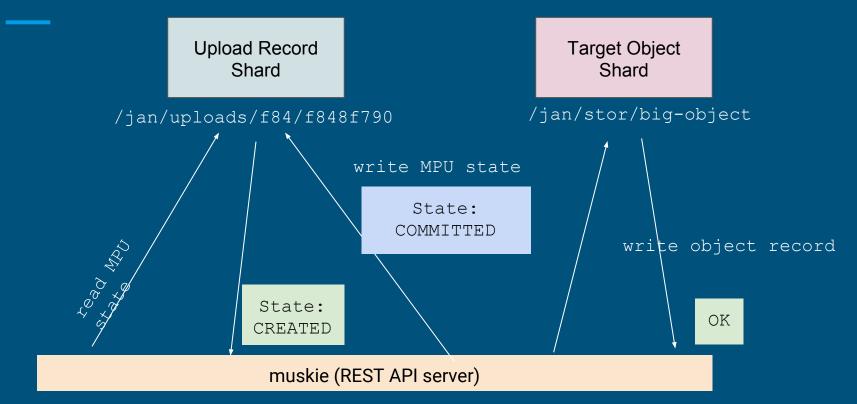
Target Object Shard

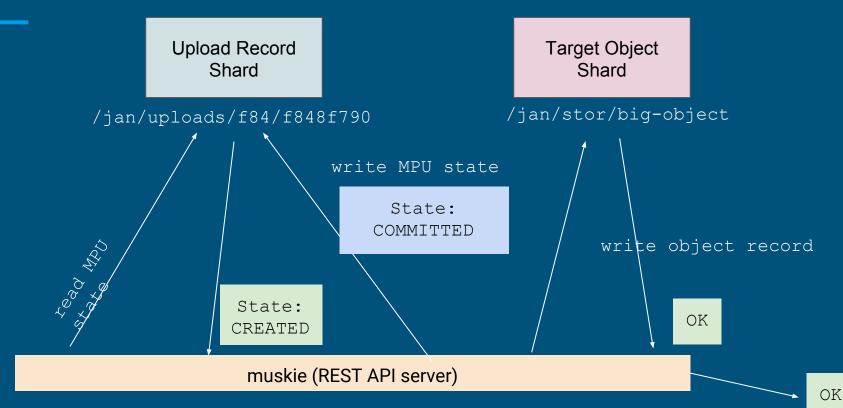
muskie (REST API server)











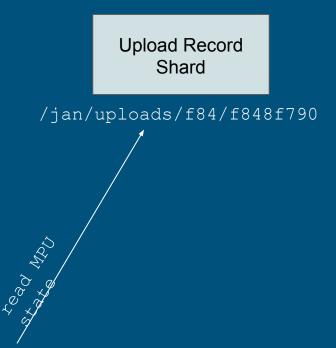


Upload Record Shard

/jan/uploads/f84/f848f790

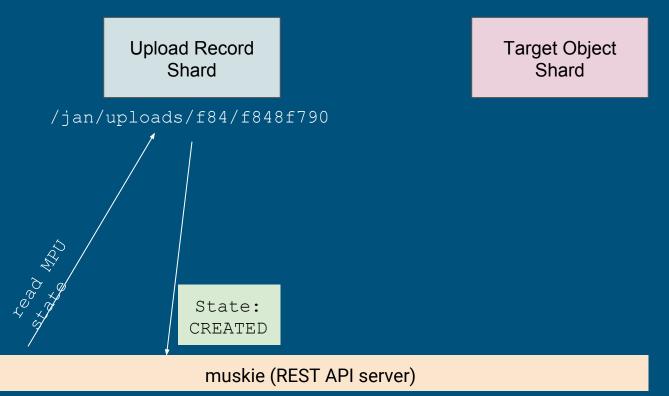
Target Object Shard

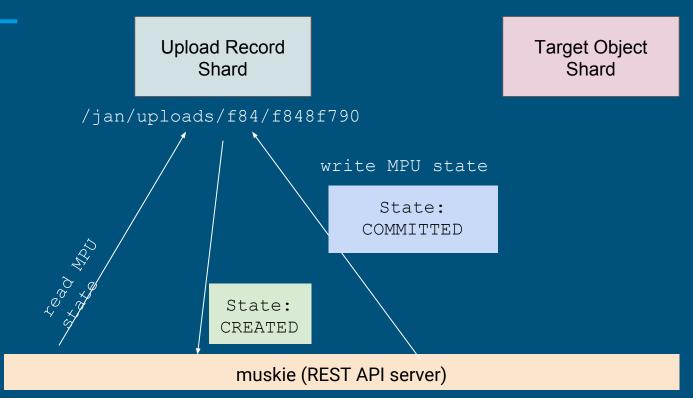
muskie (REST API server)

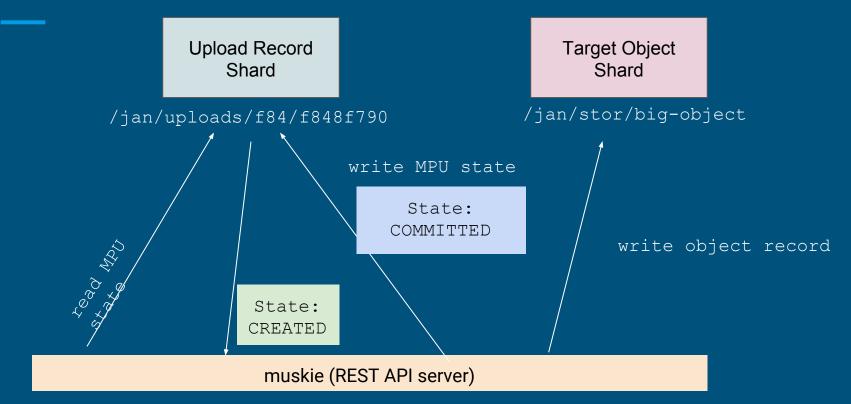


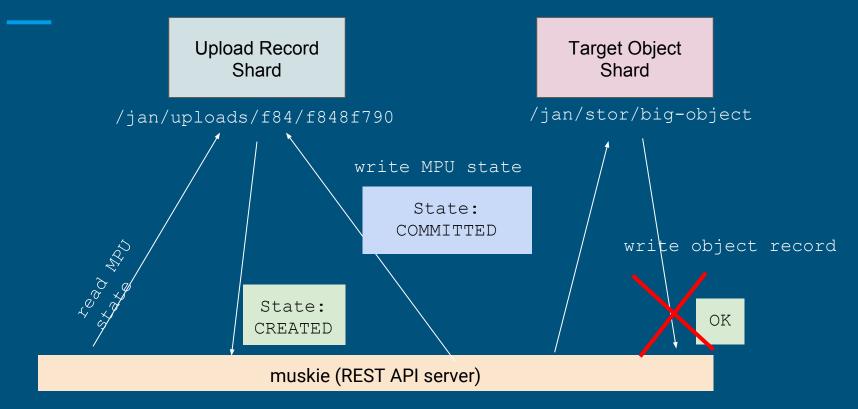
Target Object Shard

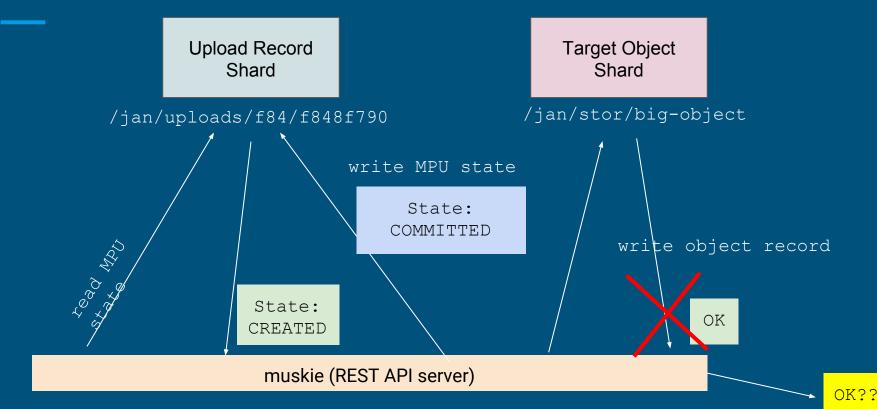
muskie (REST API server)





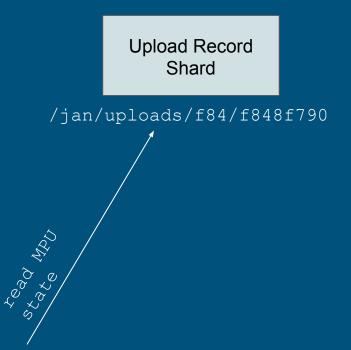






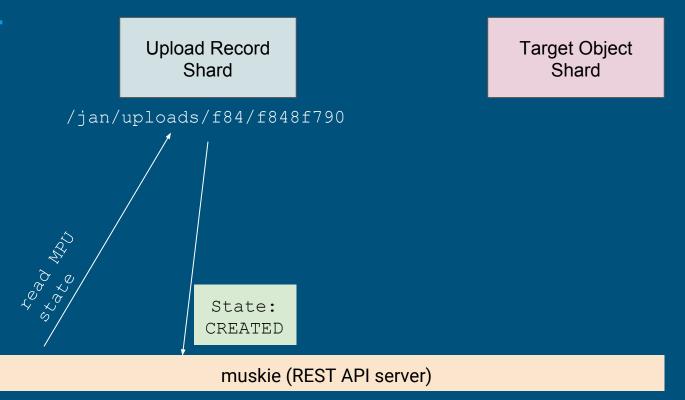
Commits: atomicity & idempotency

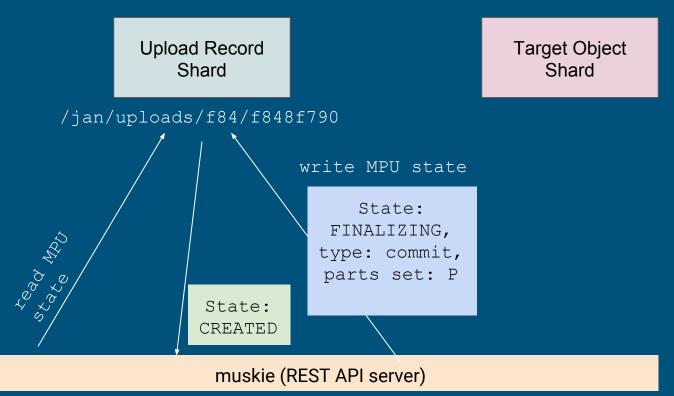
- Storing official state of MPU on parts directory metadata record is not atomic!
- Need visibility of object and state change of multipart upload to occur together, or not at all
- Solution?
 - Store official state (whether MPU is done) on the same shard as object!
 - Can still store some state on parts directory record as an optimization

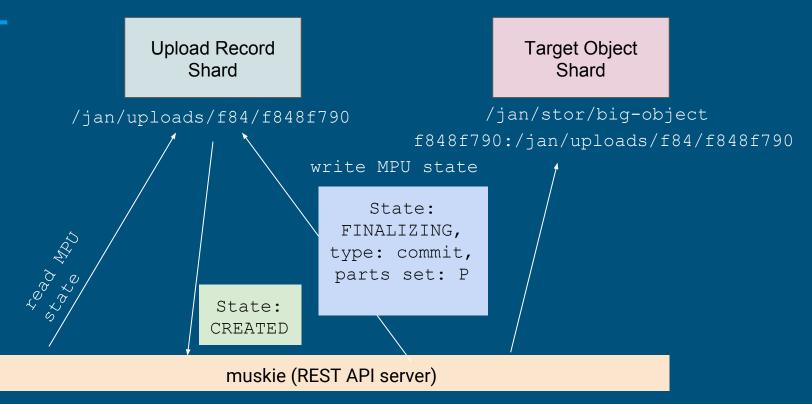


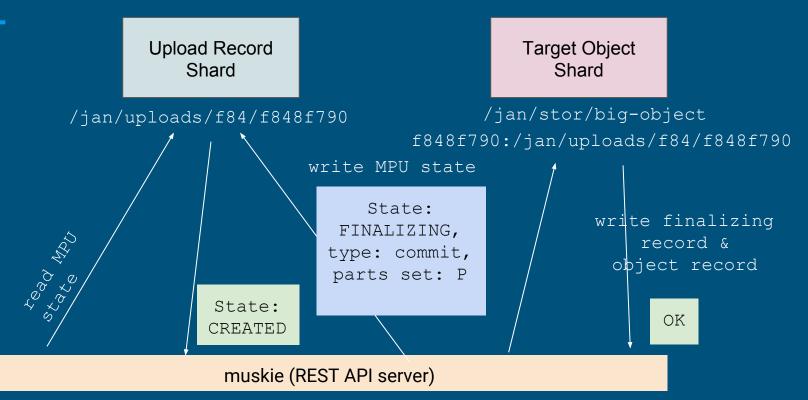
Target Object Shard

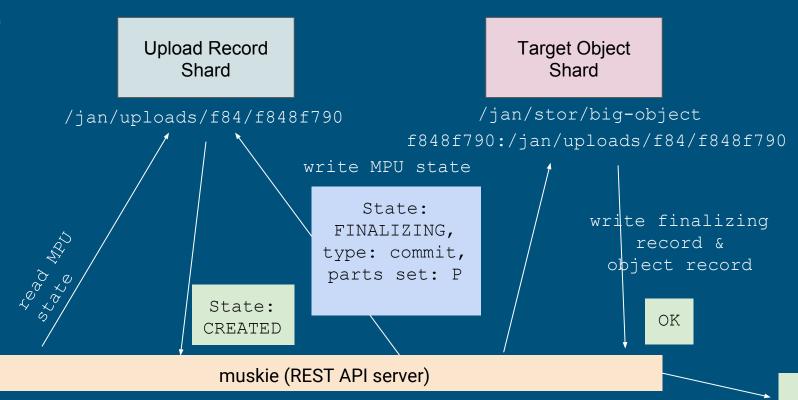
muskie (REST API server)











Upload Record Shard

Shard

Target Object

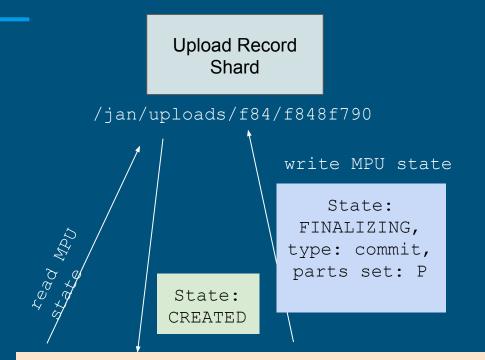
/jan/uploads/f84/f848f790



Upload Record Shard /jan/uploads/f84/f848f790 read Mpu State: CREATED

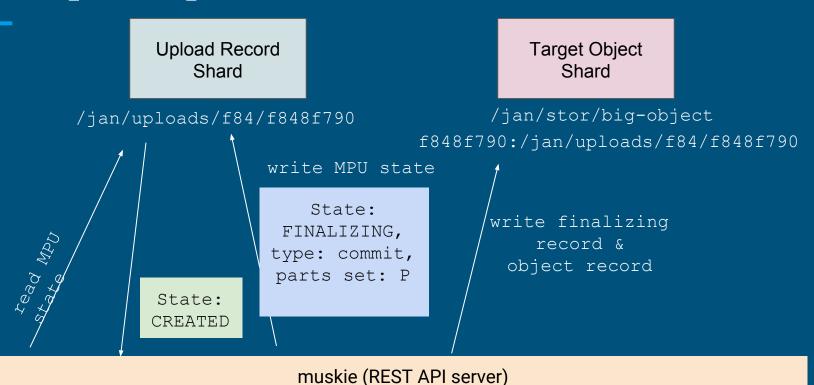
Target Object Shard

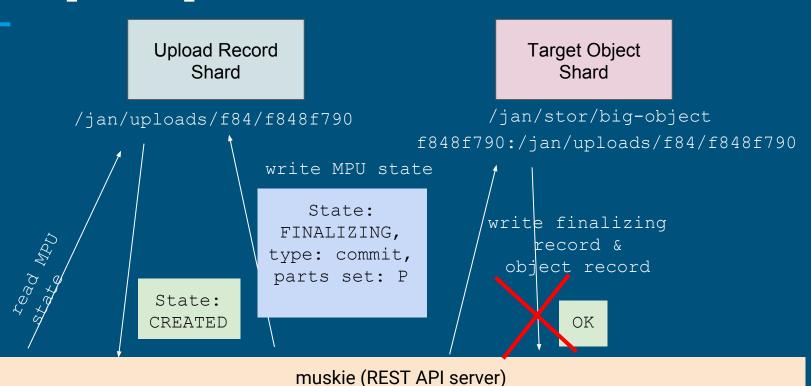
muskie (REST API server)

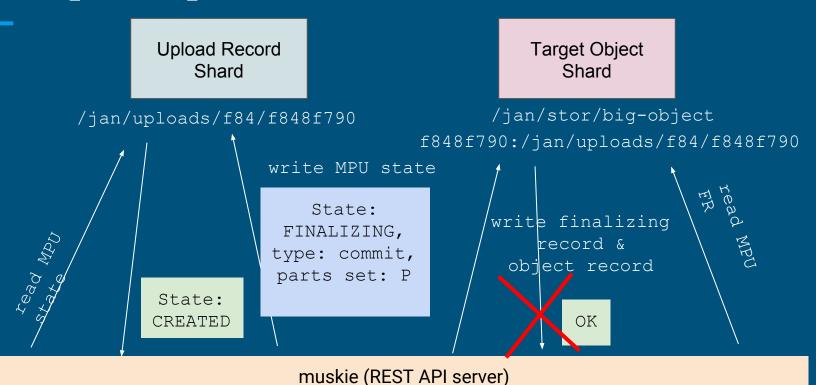


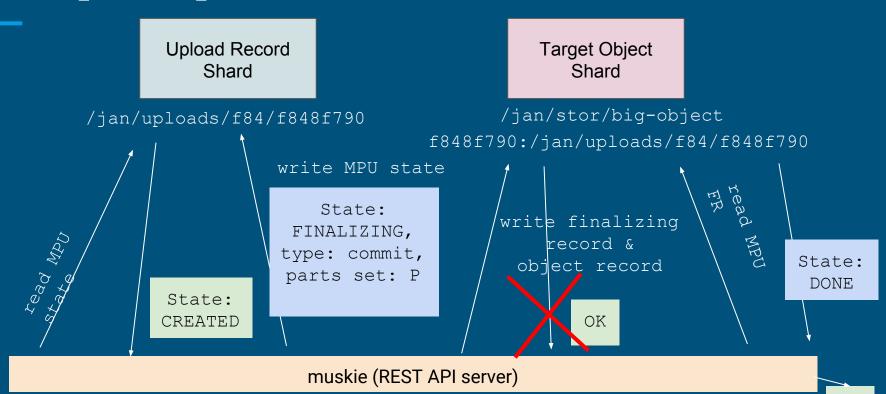
Target Object Shard

muskie (REST API server)









MPU State Machine

Commits: Final Steps

- Read upload directory record and check MPU's state.
 - If state is CREATED, verify etags of all parts.
 - If the etags are valid, update the state on the upload record to FINALIZING, type commit.
 - If the etags aren't valid, return an error.
 - If state is FINALIZING, type commit, then verify the part etags match the MD5 etag summary stored on the upload directory record. If they don't, return an error.
 - If state is FINALIZING, type abort, return an error.
- Invoke mako-finalize on the storage node set.
- Atomically insert a finalizing record and the target object record on the shard of the target object.
- Return a response indicating success to the client.

Concurrency in other MPU operations

- mpu-create: creating a multipart upload
 - Returns a unique handle for an MPU (no contention on target object path)
- mpu-get: get the state of an MPU
 - Metadata read
- upload-part: upload a part to a given MPU
 - Can overwrite parts as often as you want (consistent with Manta PUT behavior)
 - Race between updating a upload record state to "finalizing" and inserting a new part record
 - Doesn't lead to any server-size inconsistency, but likely indicates a buggy use of the API
- abort-mpu: abort an upload
 - Same concurrency protections as commit-mpu (mako-finalize is not invoked)

Revisiting Design Constraints

- Multipart upload as a feature started with its own set of constraints
 - Atomic commits/aborts
 - Sane listing of parts & uploads
- ... but was constrained to invariants of the system it was designed for
 - Immutable objects
 - Separation between metadata and storage
 - No support for cross-shard transactions
 - Composed of distributed services that can fail unexpectedly

Working with Design Constraints

- Immutable objects: mutating objects was not a possible solution
- Separation between metadata and storage: separate mechanisms of maintaining correctness of metadata layer and storage layer
- No support for cross-shard transactions: use only one shard as the final source of truth for the state of an MPU
- Composed of distributed services that can fail unexpectedly: consider atomicity & idempotency of all operations

Tradeoffs in Design

- mako-finalize: a complex, variable-latency operation added to mako (previously only a thin shim on top of nginx)
 - o Tradeoff: Variable-length latency hit for mpu-commit
 - Alternative: Copying data across the network, which would probably be slower in most cases.
- muskie (REST API service) chooses storage nodes for an MPU's target object when it is created
 - Tradeoff: Storage nodes selected may not be available when parts are uploaded or the object is committed.

Final Thoughts

- For a complete discussion of MPU design, see RFD 65
- When adding new functionality to a system, consider how it will maintain the invariants of the system (and if it doesn't, at what cost?)
- Even *legacy* systems can have new and innovative features the original authors never imagined:)
 - Can't wait for a similar presentation on Manta's tenth birthday!

Questions?